

# On the Size of Maximum $k$ -Edge-Colorable Subgraphs in Regular Graphs

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## ABSTRACT

An edge-coloring of a graph is an assignment of colors to edges of the graph, so that adjacent edges receive different colors. A subgraph  $H$  of a graph  $G$  is maximum  $k$ -edge-colorable, if  $H$  is  $k$ -edge-colorable and contains as many edges as possible. In this paper, we present some results towards the problem of finding a tight lower bound for  $\frac{|E(H)|}{|E(G)|}$ , where  $H$  is a maximum  $k$ -edge-colorable subgraph of  $G$ , and  $G$  is an arbitrary or a regular graph.

## Keywords

Edge coloring,  $k$ -edge-coloring,  $k$ -edge colorable subgraph,  $r$ -regular graph

## 1. INTRODUCTION

The classical edge coloring problem is frequently used to model real-world problems of resource allocation, such as scheduling of tasks requiring the cooperation of two processors, file transfer operations, and assignment of channels of satellite communication.

For any positive integer  $k$  there are graphs that do not admit an edge-coloring with  $k$  colors. Thus, an interesting problem is to find the maximum number of edges we can color the given  $k$  colors.

In this paper, we consider finite, undirected graphs with no loops. Graphs may contain parallel edges. For a graph  $G$ , let  $V(G)$  denote its vertex set and  $E(G)$  its edge set. The degree of a vertex  $v$  of  $G$ , denoted by  $d_G(v)$ , is the number of edges of  $G$  that are incident to  $v$ . A graph is  $r$ -regular if  $d_G(v) = r$  for all  $v \in V(G)$ . The maximum of  $d_G(v)$  over all  $v \in V(G)$  is called a maximum degree of  $G$  and is denoted by  $\Delta(G)$ . For a vertex  $v$  of  $G$  let  $\partial_G(v)$  be the set of edges of  $G$  incident to  $v$ . The girth of a graph is the length of a shortest cycle of the graph.

A matching  $M$  in  $G$  is a set of pairwise non-adjacent edges, that is, a subset of edges, where no two edges share a common vertex. A  $k$ -factor of a graph  $G$  is a spanning  $k$ -regular subgraph of  $G$ . Thus, the edge-set of a 1-factor is a perfect matching.

A  $k$ -edge coloring of a graph  $G$  is an assignment of  $k$  colors to the edges of  $G$ , so that adjacent edges receive different colors. A  $k$ -edge coloring can be thought of as

a partition  $(E_1, E_2, \dots, E_k)$  of  $E(G)$ , where  $E_i$  denotes the subset of  $E(G)$  having color  $i$ . It is not hard to see that a  $k$ -edge coloring is just a partition  $(E_1, E_2, \dots, E_k)$ , in which each subset  $E_i$  is a matching [9].

The least number  $k$  for which  $G$  has a  $k$ -edge coloring is called a chromatic index of  $G$  and is denoted by  $\chi'(G)$ . The graphs  $G$  with  $\chi'(G) = \Delta(G)$  are said to be class I, otherwise they are class II. Shannon's theorem says that for any  $G$  graph  $\Delta(G) \leq \chi'(G) \leq \left\lceil \frac{3\Delta(G)}{2} \right\rceil$  and Vizing's theorem states that for any  $G$  graph  $\Delta(G) \leq \chi'(G) \leq \Delta(G) + \mu(G)$  where the  $\mu(G)$  is the maximum multiplicity of an edge in  $G$ .

For a graph  $G$  and a positive integer  $k$ , a subgraph  $H$  of  $G$  is called maximum  $k$ -edge colorable if it is  $k$ -edge-colorable and contains as many edges as possible. The number of edges of  $H$  is denoted by  $\nu_k(G)$ . The maximum  $k$ -edge-coloring subgraph problem is the following: for a given graph  $G$  find a maximum size subgraph  $H$  of  $G$ , which is  $k$ -edge-colorable. In other words, in this problem we are looking for a  $k$ -edge-colorable subgraph containing  $\nu_k(G)$  edges.

There are several papers where the ratio  $\frac{|E(H)|}{|E(G)|}$  has been investigated. In [8], an algorithm for the problem is presented. There for each fixed value of  $k \geq 2$ , a polynomial time approximation algorithm is described, where the approximation ratios are tending to 1 as  $k$  tends to infinity. In [1], it is shown that any 2-factor of a cubic graph can be extended to a maximum 3-edge colorable subgraph. Also the authors proved that for every cubic graph  $G$

$$\nu_2(G) \geq \frac{4}{5}|V(G)| \text{ and } \nu_3(G) \geq \frac{7}{6}|V(G)|.$$

Moreover, it can be shown that

$$\nu_2(G) + \nu_3(G) \geq 2|V(G)|,$$

and

$$\nu_2(G) \leq \frac{|V(G)| + 2\nu_3(G)}{4}$$

The last equality has been investigated in [5]. There, it is shown that

$$\nu_2(G) \geq \alpha \cdot \frac{|V(G)| + 2\nu_3(G)}{4}$$

where  $\alpha = \frac{16}{17}$ , if  $G$  is a cubic graph,  $\alpha = \frac{20}{21}$  if  $G$  is a cubic graph containing a perfect matching and  $\alpha = \frac{44}{45}$  if  $G$  is a bridgeless cubic graph. There, also the improved lower bounds of  $\nu_2(G)$  and  $\nu_3(G)$  are proved when  $G$  is a claw-free bridgeless cubic graph:

$$\nu_2(G) \geq \frac{35}{36} \cdot |V(G)|, \nu_3(G) \geq \frac{43}{45} \cdot |E(G)|.$$

There are some results in [10] about maximum  $\Delta$ -edge colorable subgraph of class II graphs. There, the authors proved that every set of disjoint cycles of a class II graph with  $\Delta \geq 3$  can be extended to a maximum  $\Delta$ -edge colorable subgraph. It is also shown there that a maximum  $\Delta$ -edge colorable subgraph of a simple graph is always class I. Finally, if  $G$  is a graph with girth  $g \in \{2k, 2k+1\}$  ( $k \geq 1$ ) and  $H$  is a maximum  $\Delta$ -edge colorable subgraph of  $G$ , then

$$\frac{|E(H)|}{|E(G)|} \geq \frac{2k}{2k+1}$$

and the bound is best possible in a sense that there is an example attaining it.

Finally, let us note that the lower bounds for  $\frac{\nu_k(G)}{|V(G)|}$  in cubic graphs have been investigated in [2, 6, 11, 12, 14] when  $k = 1$ , and for regular graphs of high girth in [3]. These lower bounds have also been investigated in the case when the graphs need not be cubic [4, 7, 13].

In this paper, we prove a best-possible bound for  $\frac{|E(H)|}{|E(G)|}$  in the class of all graphs. We also investigate the same problem in the class of regular graphs and present some partial results in relation to them. Non-defined terms and concepts can be found in [9].

## 2. RESULTS

In this section, we present our main results. First, we consider the following problem in the class of all graphs.

**Problem 1.** For  $\Delta \geq 1$  and  $k = 1, \dots, \lfloor \frac{3\Delta}{2} \rfloor$  define the function  $g(\Delta, k)$  as the infimum of  $\frac{|E(H_k)|}{|E(G)|}$ , where  $G$  is any graph and  $H_k$  is a maximum  $k$ -edge-colorable subgraph of  $G$ . The infimum is taken over all graphs  $G$  of maximum degree  $\Delta$ . The problem is to determine the function  $g$ .

Our first result states:

**Theorem 1.** For  $\Delta \geq 1$  and  $k = 1, \dots, \lfloor \frac{3\Delta}{2} \rfloor$ , we have  $g(\Delta, k) = \frac{k}{\lfloor \frac{3\Delta}{2} \rfloor}$ .

*Proof.* First let us show that  $g(\Delta, k) \leq \frac{k}{\lfloor \frac{3\Delta}{2} \rfloor}$ . Consider a graph  $G$  on three vertices  $a, b$  and  $c$ , where  $a$  and  $b$  are joined with  $\lfloor \frac{\Delta}{2} \rfloor$  edges,  $a$  and  $c$  are joined with  $\lfloor \frac{\Delta}{2} \rfloor$  edges, and  $b$  and  $c$  are joined with  $\lceil \frac{\Delta}{2} \rceil$  edges. It can be easily seen that  $\frac{|E(H_k)|}{|E(G)|} = \frac{k}{\lfloor \frac{3\Delta}{2} \rfloor}$ .

In order to prove the opposite inequality, let  $G$  be any graph with maximum degree  $\Delta$ . From Shannon theorem we have  $\chi'(G) \leq \lceil \frac{3\Delta}{2} \rceil$ , thus, to prove the theorem it's enough to prove that  $\nu_k(G) \geq \frac{k}{\chi'(G)} \cdot |E(G)|$ . It will be proved using the following proposition:

**Proposition 1.** Let  $a_1 \geq \dots \geq a_n$  be any positive numbers and let  $k \leq n$ . Then the arithmetical mean of the  $a_1, \dots, a_k$  is not less than the arithmetical mean of  $a_1, \dots, a_n$ .

Consider an edge-coloring of  $G$  with  $n = \chi'(G)$  colors. Let  $a_1, \dots, a_n$  be the sizes of the color classes. We can assume that  $a_1 \geq \dots \geq a_n$ . By Proposition 1, we have

$$\frac{\nu_k(G)}{k} \geq \frac{a_1 + \dots + a_k}{k} \geq \frac{a_1 + \dots + a_n}{n} = \frac{|E(G)|}{\chi'(G)},$$

which proves the theorem.  $\square$

The second problem that we considered is the following:

**Problem 2.** For  $r \geq 3$  and  $k = 1, \dots, \lfloor \frac{3r}{2} \rfloor$  define the function  $f(r, k)$  as the infimum of  $\frac{|E(H_k)|}{|E(G)|}$ , where  $G$  is any  $r$ -regular graph and  $H_k$  is a maximum  $k$ -edge-colorable subgraph of  $G$ . The infimum is taken over all  $r$ -regular graphs  $G$ . The problem is to determine the function  $f$ .

We suspect that:

**Conjecture 1.** For  $r \geq 3$  and  $k = 1, \dots, \lfloor \frac{3r}{2} \rfloor$  one has:

$$f(r, k) = \begin{cases} \frac{2k}{3r}, & \text{if } r \text{ is even;} \\ \frac{2k(r+1)}{r(3r+1)}, & \text{if } r \text{ is odd and } k \leq \frac{3r+1}{4}; \\ \frac{2k+1}{3r}, & \text{if } r \text{ is odd and } k \geq \frac{3r+1}{4}. \end{cases}$$

Note that if  $r$  is odd and  $k = \frac{3r+1}{4}$ , then the two expressions give the same value.

Related with this conjecture, we are able to show:

**Theorem 2.** Conjecture 1 is true when  $r$  is even, or when  $r$  is odd and  $k = 1$ .

*Proof.* We start with the case when  $r$  is even. By Theorem 1, we have  $f(r, k) \geq \frac{2k}{3r}$ . On the other hand, if we consider the same example from the proof of Theorem 1 when  $\Delta = r$ , one can easily see that we get an  $r$ -regular graph, hence the converse inequality is also true. Thus,  $f(r, k) = \frac{2k}{3r}$  when  $r$  is even.

Now, let us consider the case when  $r$  is odd and  $k = 1$ . Consider the following graphs  $A_r$  and  $B_r$ . We make  $A_r$  by taking  $r$  copies of the Shannon's triangle (the graph from the proof of Theorem 1 when  $\Delta = r$ ), and taking a new vertex  $z$  and joining it to  $r$  vertices that are of degree  $r-1$ , thus we get an  $r$ -regular graph and  $B_r$  by taking two copies of the Shannon's triangle, and joining the two vertices of degree  $r-1$  with an edge. Again we get an  $r$ -regular graph. Now, it is a matter of direct verification, that if  $k \leq \frac{3r+1}{4}$  then  $A_r$  attains the bound of the conjecture, while for  $k \geq \frac{3r+1}{4}$ ,  $B_r$  attains the bound of the conjecture. Thus,  $f(r, k)$  is at most the bound predicted by the Conjecture 1, in particular, when  $k = 1$ .

Now, let us show that  $f(r, 1)$  is at least the bound of the Conjecture 1 when  $k = 1$ . Nishizeki in [11] has shown that any odd  $r$ -regular graph  $G$  contains a matching of size at least  $\lceil \frac{(r^2-r-1)|V|-(r-1)}{r(3r-5)} \rceil$ . Now it can be shown that this expression is at least  $\frac{r+1}{3r+1} \cdot |V(G)| = \frac{2(r+1)}{r(3r+1)} \cdot |E(G)|$ . Thus,  $f(r, 1) \geq \frac{2(r+1)}{r(3r+1)}$ .  $\square$

Finally, it turns out that there is a certain dependence among different values of  $k$  in the Conjecture 1. More precisely, we are able to show that

**Theorem 3.** *If Conjecture 1 is true when  $r$  is odd and  $k = \lfloor \frac{3r+1}{4} \rfloor$ , then it is true when  $r$  is odd and  $k = 1, \dots, \lfloor \frac{3r+1}{4} \rfloor$ .*

*Proof.* Assume that  $r$  is odd. Observe that for  $i = 1, \dots, k$ , where  $k = \lfloor \frac{3r+1}{4} \rfloor$ , we have  $\nu_i(G) \geq \frac{i}{k} \cdot \nu_k(G)$ . Hence, if Conjecture 1 is true when  $k = \lfloor \frac{3r+1}{4} \rfloor$  then it is true for  $i = 1, \dots, k$ .  $\square$

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