# An Upper Bound on the Edge-Chromatic Sum of Fibonacci Cubes 

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#### Abstract

Fibonacci cube is an isometric subgraph of the $n$-dimensional cube. A proper edge-coloring of a graph $G$ is a mapping $\alpha: E(G) \longrightarrow \mathbb{N}$ such that $\alpha(e) \neq \alpha\left(e^{\prime}\right)$ for every pair of adjacent edges $e$ and $e^{\prime}$ in $G$. The edge-chromatic sum of a graph $G$ is the minimum sum of all colors in the graph among all its proper edge-colorings. This paper provides an upper bound on the edge-chromatic sum of Fibonacci cubes.


Keywords- Edge-chromatic sum, Fibonacci cubes, sum edgecoloring.

## I. Introduction

Let $B=\{0,1\}$ and for $n \geq 1$ set $\mathcal{B}_{n}=\left\{b_{1} b_{2} \ldots b_{n} \mid b_{i} \in\right.$ $B, 1 \leq i \leq n\}$. The $n$-cube $Q_{n}$ graph is the graph defined on the vertex set $\mathcal{B}_{n}$, vertices $b_{1} b_{2} \ldots b_{n}$ and $b_{1}^{\prime} b_{2}^{\prime} \ldots b_{n}^{\prime}$ being adjacent if $b_{i} \neq b_{i}^{\prime}$ holds for exactly one $i \in\{1,2, \ldots, n\}$.

Fibonacci cubes are introduced as follows: for $n \geq 1$, let $\mathcal{F}_{n}=\left\{b_{1} b_{2} \ldots b_{n} \in \mathcal{B}_{n} \mid b_{i} \cdot b_{i+1}=0,1 \leq i \leq n-1\right\}$. The Fibonacci cube $\Gamma_{n}$ is the subgraph of $Q_{n}$ induced by the vertex set $\mathcal{F}_{n}[1]$.

A proper vertex-coloring of a graph $G$ is a mapping $\alpha: V(G) \rightarrow \mathbb{N}$ such that $\alpha(u) \neq \alpha(v)$ for every $u v \in E(G)$. In that case $\alpha(v)$ is called the color of the vertex $v$. The vertex-chromatic sum $\Sigma(G)$ of a graph $G$ is the minimum sum of colors of all vertices among all proper vertex-colorings of $G$. The concept of vertex-chromatic sum was introduced by Kubicka [2] and Supowit [3]. The problem of finding the vertex-chromatic sum is shown to be NP-complete in general and polynomial time solvable for trees [4]. Jansen [5] gave a dynamic programming algorithm for partial $k$-trees. In papers [6], [7], [8], [9], [10], some approximation algorithms were given for various classes of graphs. Some bounds for the vertex-chromatic sum of a graph were given in [11].

A proper edge-coloring of a graph $G$ is a mapping $\alpha$ : $E(G) \rightarrow \mathbb{N}$ such that $\alpha(e) \neq \alpha\left(e^{\prime}\right)$ for every adjacent $e$ and $e^{\prime}$. In that case $\alpha(e)$ is called a color of the edge $e$. Similar to the vertex-chromatic sum of graphs, in [6], [12], and [13], edge-chromatic sum of graphs was introduced. Namely, the edge-chromatic sum $\Sigma^{\prime}(G)$ of a graph $G$ is the minimum sum of all colors in the graph among all its proper edge-colorings. In [6], Bar-Noy et al. proved that the problem of finding the edge-chromatic sum is NP-hard for multigraphs. Later, in [12], it was shown that the problem is NP-complete for bipartite graphs with maximum degree 3. Petrosyan and Kamalian [14]
proved that the problem is NP-complete for even more specific class of graphs from the latter and found an $\frac{11}{8}$-approximation algorithm for $r$-regular graphs. In [15], Salavatipour proved that determining the edge-chromatic sum is NP-complete for $r$-regular graphs with $r \geq 3$. The problem can be solved in polynomial time for trees [12].

The terms and concepts that we do not define can be found in [16].

In the present paper, we obtain an upper bound on the edgechromatic sum of Fibonacci cubes.

## II. Main result

Theorem 1. For any $n \in \mathbb{N}$, we have

$$
\begin{gathered}
\Sigma^{\prime}\left(\Gamma_{n}\right) \leq \\
\leq\left(\frac{5+3 \sqrt{5}}{100} n^{2}+\frac{11+9 \sqrt{5}}{100} n+\frac{6 \sqrt{5}}{125}\right)\left(\frac{1+\sqrt{5}}{2}\right)^{n}+ \\
+\left(\frac{5-3 \sqrt{5}}{100} n^{2}+\frac{11-9 \sqrt{5}}{100} n-\frac{6 \sqrt{5}}{125}\right)\left(\frac{1-\sqrt{5}}{2}\right)^{n} .
\end{gathered}
$$

Proof. We will construct a corresponding proper edgecoloring $\alpha_{n}$ for each $\Gamma_{n}$ by induction on $n$. We denote by $\Sigma^{\prime}\left(\alpha_{n}\right)$ the sum of colors of all edges for the coloring $\alpha_{n}$. From the relation $\Sigma^{\prime}\left(\Gamma_{n}\right) \leq \Sigma^{\prime}\left(\alpha_{n}\right)$, which implies from the definition of the edge-chromatic sum, we will derive the result.

It is easy to construct $\alpha_{1}$ and $\alpha_{2}$ separately, so they have, respectively, 1 and 3 as their sums.

Now let us construct the coloring $\alpha_{n}$ for $n \geq 3$ assuming that we have already constructed all $\alpha_{k}$-s for $1 \leq k<n$. It is known that it is possible to decompose $\Gamma_{n}$ into two subgraphs $\Gamma_{n-1}$ and $\Gamma_{n-2}$ in such a way that $V\left(\Gamma_{n}\right)=V\left(\Gamma_{n-2}\right) \cup$ $V\left(\Gamma_{n-1}\right)$ and $E\left(\Gamma_{n}\right)=E\left(\Gamma_{n-2}\right) \cup E\left(\Gamma_{n-1}\right) \cup M$, where $M$ is a matching of $2\left|V\left(\Gamma_{n-2}\right)\right|$ vertices [1]. Let us color the edges of the matching with the color 1 . For the remaining edges let us use the corresponding colors in the colorings $\alpha_{n-2}$ and $\alpha_{n-1}$, and color the edge $e$ with $\alpha_{n-2}(e)+1$ if $e \in E\left(\Gamma_{n-2}\right)$ and $\alpha_{n-1}(e)+1$ if $e \in E\left(\Gamma_{n-1}\right)$.

Clearly, we constructed a proper edge-coloring. Moreover, $\Sigma^{\prime}\left(\alpha_{n}\right)=\left|E\left(\Gamma_{n}\right)\right|+\Sigma^{\prime}\left(\alpha_{n-1}\right)+\Sigma^{\prime}\left(\alpha_{n-2}\right)$. By [1], we have that $\left|E\left(\Gamma_{n}\right)\right|=\frac{n F_{n+1}+2(n+1) F_{n}}{5}$, where $F_{n}$ is the $n$-th Fibonacci number.

If we denote $\Sigma^{\prime}\left(\alpha_{n}\right)$ by $a(n)$, then to obtain the required inequality, it is necessary to solve the 2 nd order nonhomogeneous recurrence relation $a(n)=a(n-1)+a(n-2)+$ $\frac{n F_{n+1}+2(n+1) F_{n}}{5}$ with initial conditions $a(1)=1$ and $a(2)=$ 3. To do that, we represent $a(n)$ as the sum $a_{h}(n)+a_{p}(n)$, where $a_{h}(n)$ is the solution of the associated homogeneous recurrence relation, and $a_{p}(n)$ is the particular solution. The characteristic equation of the homogeneous relation is $\lambda^{2}-\lambda-$ 1 , roots of which are $\lambda=\frac{1+\sqrt{5}}{2}$ and $\lambda=\frac{1-\sqrt{5}}{2}$. Therefore, $a_{h}(n)=c_{1}\left(\frac{1+\sqrt{5}}{2}\right)^{n}+c_{2}\left(\frac{1-\sqrt{5}}{2}\right)^{n}$. As for the particular solution, considering the formula of the common term of the Fibonacci sequence, we get that $a_{p}(n)$ has the following form: $\left(A n^{2}+B n+C\right)\left(\frac{1+\sqrt{5}}{2}\right)^{n}+\left(D n^{2}+E n+F\right)\left(\frac{1-\sqrt{5}}{2}\right)^{n}$. Thus, we obtained the forms of $a_{h}(n)$ and $a_{p}(n)$, and to get the formula of $a(n)$ we need to put those results in the recurrence relation, and obtain the values of the coefficients using the initial conditions.

The proper edge-colorings $\alpha_{3}$ and $\alpha_{4}$ are illustrated in Figure 1.


Figure 1

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